

Contents

Introduction	5
Chapter 1: symbols	6
What is the symbol	6
The structure symbol (SYMBOL).....	6
Chapter 2: The format of binary blocks SpartaDOS	8
The types of blocks of binary	8
Block booking memory	8
Relokowalny block binary	9
Block fixupów	9
The lists of required symbols	10
The definition of a symbol	11
Chapter 3: memory management	12
The use of memory by SpartaDOS	12 ..
Recognizing the configuration memory	14
The list of free memory banks (T_).....	16
The allocation of main memory and an additional (malloc).....	18
The allocation of memory banks	19
Access to the bank system	20
Access to other banks enlargement	22
Chapter 4: Handling errors	24
The standard procedure for handling error (U_FAIL).....	24
The creation of traps (U_SFALL).....	24
Removing the trap (U_XFAIL).....	25
Error message (U_ERROR).....	25
An example of the use of traps	25
Chapter 5: Command-line treatment	27
LBUF and buffer index BUFOFF	27
Buffer COMFNAM	27
Reading the command line	28
The number (U_GETNUM).....	28
The dwustanowy: ON and OFF (U_GONOFF).....	29
Options (U_SLASH).....	29
Keyword (U_TOKEN).....	30
Device Name (U_GETPAR / U_GEFINA).....	31
The specification of the file (U_GETPAR / U_GEFINA).....	32
Specification directory (U_GETPAR / U_GEPATH).....	32
The specification and file attributes (U_GETATR).....	33
Specification file with the default Mask (U_FSPEC).....	34
Combinations of different types of parameters	34
Chapter 6: processing the file name	36
Converting from 8 +3 format for internal (PRO_NAME).....	36
Converting from an internal format for 8 +3 (U_EXPAND).....	36

What is the environment variable	37
Read variable by its name (getenv).....	37
Read by a variable number (NUMENV).....	37
Record and erase the variables (putenv).....	38
Chapter 8: read and write files	39
Opening the files (fopen).....	39
Closing files (fclose / FCLOSEAL).....	40
Read and write individual bytes (FGETC / FPUTC).....	41
Read and write records (fgets / FPUTS).....	41 ..
Record formatted text file (FPRINTF).....	42
Read and write blocks of binary (fread / fwrite).....	42
Read the length of the file (FILELENG).....	43
Changing positions in the file (FTELL / FSEEK).....	43
Chapter 9: console, entry and exit	45
Record of individual characters on the screen (PUTC).....	45
Record record on the screen (PUTS).....	45
Record formatted text on the screen (printf).....	45
Read a single character to the console (GETC).....	49
Read the record from the console (GETS).....	49
Redirection of I / O (DIVIO / XDIVIO).....	50
Vector output (PUT_V / VPRINTF).....	51
Chapter 10: catalogs	52
Search files (FFIRST / FNEXT).....	52 ..
Read the catalog (FDOPEN / FDGETC / FDCLOSE).....	52
Chapter 11: loading of binary	55
Load the program into memory (U_LOAD).....	55
To remove the program from memory (U_UNLOAD).....	55
Chapter 12: File management functions	56
Change the file name (RENAME).....	56
Deleting a file (REMOVE).....	56
Removing directory (RMDIR).....	56
The creation of a directory (MKDIR).....	57
Changing the attributes (chmod).....	57
Make sure you download the BOOT (SETBOOT).....	58
Chapter 13: Other features disc	59
Formatting the drive (FORMAT).....	59
Record fresh directory (builddir).....	59
Can not read the disk parameters (GETDFREE).....	60
Change the current directory (CHDIR).....	61
Read the current directory (getcwd).....	61
Chapter 14: ancillary functions	62
32-bit multiplication and division (MUL_32/DIV_32).....	62
Replacing the small letters on the large (ToUpper).....	62
Checking the directory separator (CKSPEC).....	62
Chapter 15: procedures for initiating	63

Initiating overlays after RESET (S_ADDIZ).....	63
Exit to DOS-u (_DOS).....	63
Warm restart SpartaDOS (_INITZ).....	63
Chapter 16: the manipulation of symbols list	65
Searching the list of symbols (S_LOOKUP).....	65
The addition symbol (S_ADD).....	66
Stripping (S_CLEAR).....	67
Chapter 17: other symbols library	68
Index procedures and system variables69

Introduction

This manual is intended for coders to

'd like to write applications and system for SpartaDOS X.

The authors assume that the reader has the idea of writing programs

Atari, and is an advanced user SpartaDOS X, to

What message contained in the "Disk Operating System

SpartaDOS X User's Guide "and supplement" SpartaDOS

V. X 4.40 "omit here as self-evident. In addition, it is assumed that

concepts such as "PORTB" does not require an explanation.

Listings are written in Assembler Pseudocode the syntax is consistent with

MAE assembler. It is most similar to the syntax of assembler

MAC/65, only the details differ from the syntax used by

crossassembler MADS. A programmer wants to adapt his own hand

them to your favorite assembler.

We wish you a pleasant reading

DLT Ltd.

Chapter 1: Symbols

What is the symbol

Manual programming SpartaDOS X is a good start with

explanation of the concept of *a symbol* with which the reader will meet

while reading.

The symbol in SpartaDOS X is a kind of rich-pointer.

It contains information on where the computer's memory is the

object: variable, table or procedure. One symbol represents

one such object. All symbols are about hundreds,

are linked in the list. It also guarantees access to the most important - from the point of

programmer's view - the procedures and internal structures SpartaDOS X

and the drivers and overlays, since the resident may,

naturally, to add to the list of their own symbols.

The main advantage of such a solution is that the procedures

system, so that they are indicated by symbols, are not

permanently assigned to specific addresses. Addresses can be

vary from version to version of DOS, and yet the programs will work.

What's more, thanks to define a new symbol of the same name,

as one of the existing overlay can easily take over mediation

between the program and procedure system, which he produces.

The structure symbol (Symbol)

The exact description of the symbol is a not-for-nothing needed, because the programs do not benefit from them as to require a Programmer its knowledge. However, this information can at least help satisfy curiosity readers.

A single symbol occupies 13 bytes, consisting of him as follows:

- + \$ 00 - \$ 01 rate for the next symbol (2 bytes)
- + \$ 02 - \$ 09: the name of the symbol (8 ASCII characters)
- + \$ 0A: byte control
- + \$ 0B-\$ 0C: address indicated by the symbol (2 bytes)

Symbol name meet the same conditions as the file name: is it to eight characters ASCII, but they are generally large and sign letters of the alphabet underscores. When a name is no less than eight characters, is completed spaces.

Byte contains **control** bits in the two oldest index memory indicated by the symbol: this is 0 for main memory and 2 notes. Six bits are known. **PID (Program Identifier):** the number of the program, which consists of a symbol. Zero represents the library system, which will still question.

The functions of libraries allow access to the list of symbols are described in a subsequent chapters.

NOTE: at the time of the launching of the X command list of symbols becomes unavailable.

Chapter 2: The format of binary blocks SpartaDOS

The types of blocks of binary

Atari DOS knows only one type of binary block, it is sześciobajtowym segment with a header starting (optional indeed), from FFFF \$ signature. This format is known, therefore his description is omitted.

SpartaDOS distinguishes a total of seven types of blocks of binary.

One of them is, naturally, above-mentioned segment Atari DOS.

Nierelokowalny SpartaDOS segment is the same structure, only header signature is different: instead of \$ FFFA \$ FFFF. Changing signatures is aimed at preventing such binary read by other DOS-y, because this format is intended for system files (drivers, etc.) SpartaDOS.

The other five types are:

- 1) memory block booking
- 2) relokowalny block binary
- 3) block update addresses inside the program (the so-called *fixupy*)
- 4) lists the symbols required by the program
- 5) lists the symbols defined by the program

Block booking memory

Block booking memory causes, as taught by the same name,

set aside by the loader indicated quantity indicated memory. Block

This is due to a header from a population of only eight bytes:

+ \$ 00 - \$ 01: \$ signature FFFE

+ \$ 02: sequence block

+ \$ 03: byte control with a value of \$ 80 index added memory

+ \$ 04 - \$ 05: the beginning of the transfer terms

+ \$ 06 - \$ 07: the number of bytes to book

The memory is reserved on the index of free memory

(ie over MEMLO in the main memory). Type of memory, the main or

Secondary is indicated by an index that is in bajcie

control. The indexes more memory is written in chapter

"The economy of memory."

Relokowalny block binary

Relokowalny block binary is heading in principle the same as memory block booking.

+ \$ 00 - \$ 01: \$ signature FFFE

+ \$ 02: sequence block

+ \$ 03: byte control with a value of \$ 00 index added memory

+ \$ 04 - \$ 05: the beginning of the transfer terms

+ \$ 06 - \$ 07: the number of bytes to load

Both types of units differ only two things: first, byte

control in the header block relokowalnego is zgaszony bit 7; is

signal to the loader, that are headed for the data to load.

The second thing is precisely this fact. The binary data are loaded into the address

indicated by the rate of free memory. Type of memory, the main or

Secondary is indicated by the index recorded in bajcie control.

Block fixupów

Block update addresses internal, so called. *Fixupów*, block

relokowalnego is pięciobajtowy headline:

+ \$ 00 - \$ 01: \$ signature FFFD

+ \$ 02: the number of the next block, which fixupowanie

+ \$ 03 - \$ 04: two bytes reserved

The next bytes to transfer within the block, first in relation

to address the loading of the block, each of the next in relation to

the previous one. In other words, such a byte means "increase the address and so on

bytes and so do the fixup. " Fixupowanie lies in the fact that the address

charging unit is added to each word dwubajtowego

below the address indicated by the byte transfer.

Byte transfer shows always double-byte word - and so

Program relokowalny MAY NOT include the value of any of its

internal addresses divided into junior and senior byte located

separately. Structures such as:

lda # <address

ldx #> address

whether the plates together separately younger and older bytes indicators

are excluded.

Such is the importance of each value of less than \$ FC. The values of \$ FC to \$ FF is an additional control codes:

\$ FF: increase by 250 bytes of address (anything beyond this is not done)

\$ FE: change the number of block fixupowanego the value seq. byte

\$ FD: change the beginning of the block to the address contained in the following. word, and fixup

\$ FC: mark the end of the block

The lists of required symbols

This is a list of symbols, which must be defined in the system, to loaded the program can work. The absence of any cause interruption No charge and error 154 (Loader: The symbol not defined).

+ \$ 00 - \$ 01: \$ signature FFFB

+ \$ 02 - \$ 09: the name or symbol (8 characters)

+ \$ 0A-\$ 0B: two bytes reserved

Next followed bytes transfers and control in the same way as in the block fixupów. Address indicated by the symbol is added to dwubajtowego words below the address indicated by the byte transfer.

The definition of a symbol

Block definition symbol to add the new loader

symbol of the global list of symbols. Structure:

\$ 00 - \$ 01: \$ signature FFFC

\$ 02: the number of the block, where the symbol is defined

\$ 03 - \$ 04: Transfer

\$ 05 - \$ 0C: the name of the symbol

The address pointed to by a new symbol is the address block loading number listed in bajcie \$ 02, plus transfer of bytes \$ 03 - \$ 04

Chapter 3: memory management

The use of memory by SpartaDOS

SpartaDOS X distinguishes the following types of memory:

1) Main memory: basic 48k of RAM to MEMLO

MEMTOP-u

2) the extension of memory: 14k additional RAM ("under the ROM-height", Areas \$ C000-\$ CFFF and \$ D800-\$ FFFF) in computers 800XL and 65XE, or additional RAM (Area \$ 4000 - \$ 7FFF) for up to 1 MB in 130XE computers, or up to 2 MB in the Atari 800 computers In the absence of this memory instead of the allocated memory is Home.

3) own module ROM: 128k ROM in the area of \$ A000-\$ BFFF.

System for their own needs does the memory as follows:

1) kernel main memory, from \$ 0,700 to MEMLO.

2) drivers additional system memory

3) The library system in the module ROM

4) The device CAR: *ibid*

Already mentioned earlier *system library* contains a collection of the intermediary between the user and load DOS, and one of the software application - namely *SpartaDOS Formatter*. The library occupies two banks (after 8k) module ROM SpartaDOS X 4.20, and 3 banks in SpartaDOS X 4.40. Action and use of the library is devoted to the main part of this manual.

Leading module library (the bank No. 1) is normally present in Under \$ A000-\$ BFFF, but it can be unplugged and replaced RAM, if the program is run a command X.

Usually this happens in the case of ordinary binaries intended Atari DOS-in (those with a header \$ FFFF).

The allocation of RAM for the various elements of the system The following depending on whether it is available, and from what he chose you in your CONFIG.SYS. Here are four possible combinations:

1) the use of main memory: use none in the CONFIG.SYS.

All components of DOS are loaded in the main memory ranging from \$ 0700. address This configuration is chosen automatically on computers Atari 800 with no more than 48k of RAM.

2) the use of memory "under the ROM-em": USE OSRAM in CONFIG.SYS. Kernel memory addresses from \$ 0,700 to MEMLO, other DOS-in components are loaded into memory in the area of \$ E400-\$ FFFF, area \$ D800-\$ DFFF occupied on data (from version 4.40 -- was previously vacant), the area of \$ C000-\$ CFFF remains slow for the rest of the system is assigned the main memory. This configuration is chosen automatically on computers Atari 800XL and Atari 65XE with no more than 64k of RAM.

3) the use of memory "under the ROM-em": USE OSRAM / DEVICE SPARTA OSRAM in CONFIG.SYS. As above, except that the memory of \$ C000 - \$ CFFF utilization is at the DOS buffers respectively in reducing occupancy main memory.

4) the use of memory bankowanej: kernel is loaded in the area from \$ 0,700 to address MEMLO, additional components of DOS, data and buffers occupy an additional memory bank located in Under \$ 4000 - \$ 7FFF. This configuration is automatically selected, Atari 800 when the computer is at all equipped with such an extension

(Axlon type, up to 2 MB of RAM banks), or when the computer XL / XE is More than 128k of RAM.

Additional memory is distinguished by DOS on the bank system where the drivers reside, above all, the procedure SPARTA.SYS, and all the rest. The procedures in DOS ensure easy access only to the bank system. This is justified by the fact that memory allotted for the system may be in the main area

RAM, the ROM-in memory or in bankowanej, connecting adequate memory system takes so himself. In contrast, "the whole the rest "to extend the type 130XE (or Axlon). SpartaDOS offers here some support - as below - but the switch banks must do it on their own by interfering in the relevant records I/O.

Recognizing the configuration memory

Even programs not in principle for SpartaDOS X may be interested in taking the current configuration of the memory DOS, and especially that expanded memory banks are occupied by it. Part of this information is contained in an array COMTAB DOSVEC indicated by the vector (\$ 0A-\$ 0B), and by the symbol COMTAB.

Caution should refrain from treatment value DOSVEC as a constant. The address pointed to by this vector is different in different SpartaDOS versions (and even different versions SpartaDOS X) and not is no guarantee that it will not change in the future. Solids are only shifts (offsets) for the specified address, as shown below.

COMTAB + \$ 1D (NBANKS) contains a number of free banks additional memory type or 130XE (for the Atari 800) Axlon. Where is here is zero, this means that additional memory or does not exist, or is

fully occupied by components DOS, ie, drivers, etc. ramdyski

COMTAB + \$ 1E (BANKFLG): If there is \$ FF, the system loaded is the memory bankowanej (USE BANKED).

COMTAB + \$ 1F (OSRMFLG): If there is \$ FF, the system loaded to the memory "under the ROM-height" (USE OSRAM). Memory bank, the if any, may be used as ramdysk and allocated drivers such as CON64.SYS and CON80.SYS.

In addition, the developer may want to check out the type of enlargement Memory: When the **COMTAB + \$ 1B (_800FLG)** is \$ FF, we have to dealing with a computer Atari 800 and a Axlon. In Otherwise (if the _800FLG is zero) is a computer XL / XE with a type of 130XE.

The calculation of how much the system in general is an additional memory (used, or not) is possible on the basis of so-called masks PORTB (PBMASK). It can be found at COMTAB + \$ 1C, and includes one in all the bits, which makes set-up in the PORTB Switching memory banks in the area of \$ 4000 - \$ 7FFF. What counts is also bit 4 of the port, so for example, simply 130XE mask is set to \$ 1C, namely% 00011100. Except bit 4 are set here two bits, which means four additional banks of memory, or 64k.

By analogy with the extension to the 320k type of RAM Compy Shop mask has a value of \$ DC, which is 11011100.% Except bit 4 are here 4 bits set. This means 16 additional memory banks.

Information, which specifically memory bank notes is "The system" is contained at T_ + \$ 06 or \$ COMTAB-0150

1)

1 The only appropriate method of access to the array reference T_ is for

T_ symbol means. However, due to programs
Not for SpartaDOS X, and wishing to take advantage of these instructions
data, we - as an exception - an alternative method of calculating its
address, namely on the basis of a shift in relation to COMTAB
(ie, in relation to the value of the vector DOSVEC). Should not be there with
conclude that each of the symbols is achievable this way, rather the contrary,
the location of all other defined objects symbolically

For use none is \$ FF, and for the use and BANKED USE
OSRAM there is value to be put into the registry
PORTB, to connect the bank, in which resides the main driver
SpartaDOS (ie SPARTA.SYS).

The list of free memory banks (T_)

Obtaining information reverse, ie not that banks are busy, but
rather, which are free, is slightly more complicated. Nevertheless,
it's a lot more useful, because they are occupied by the bank
SPARTA.SYS need not be the only one in which they reside drivers DOS
u. By now skip the question ramdisków, SpartaDOS X from version 4.41 is
Two Drivers, CON64.SYS and CON80.SYS, which allocates each
additional memory banks. Overwrite them anything at the moment
when they are activated, of course, result in suspension of computer
sooner or later.

The main entrance to calculate the list of free memory banks
their number is shown by NBANKS (COMTAB + \$ 1D). For the calculation
necessary data are also included in the array T_ and the register PBMASK
(COMTAB + \$ 1C). The procedure produces a list of banks is free
as follows:

```
; FREELIST
```

```
sav
```

```
= $ 80
```

```
savy
```

```
= $ 81
```

```
temp
```

```
= $ 82
```

```
index = $ 83
```

```
lda COMTAB + $ 1d
```

```
Solid sav
```

```
LDY # $ 00
```

```
Jan savy
```

(variables, arrays, procedures) is nohow to address and may COMTAB
be any change in future versions of SpartaDOS development.

```
January index
```

```
loop
```

```
LDY sav
```

```
dey
```

```
Jan sav
```

```

bmi exit
yeah
and # $ 03
asl
asl
Solid temp
yeah
LSR
LSR
tax
lda T_ + $ 08, x
now temp
eor PORTB
and COMTAB + $ 1C
eor PORTB
PHA
iny
yeah
LDY index
Solid axlon, and
pla
Solid list, and
inc index
bne loop
exit
rts
axlon. SB 128
letter
. SB 128

```

In the place labeled "list" procedure will list PORTB registry values corresponding to free memory banks, while under the label "axlon" is corresponding with her list Axlon banks enlargement. Number of entries will be equal to the number of free banks shown by NBANKS (COMTAB + \$ 1D). Banks in accordance with the 130XE (or first 64k extension) will be on this list featured as the last one. This is in line with the overall logic of this allocation memory, according to the first-come, banks far, so that

for example on a computer equipped with 192k of RAM is SpartaDOS ulokował in the area over the basic 128k, and banks remained free 130XE programs for wanting to possibly benefit from it.

The allocation of main memory and an additional (malloc)

There are few roads, which installs the program may be rezydentnie SpartaDOS tell, that took for himself a piece of RAM.

Regular programs in binary format Atari DOS-u (with a header \$ FFFF) may, in order to raise the rate MEMLO. After paying controls for DOS in memory of information (contained in the table H_FENCE) will be created on this basis.

The second way is available for relocatable binaries SpartaDOS,

which are always loaded in the place indicated vector MEMLO -- or, more precisely, in the place indicated by the first word plate H_FENCE. Automatic allocation of the area occupied by the program is obtained by inserting a \$ FF to the indicated variable symbol INSTALL and return control to the DOS-u. Before you start scheme, resetting the variable INSTALL, so the normal way "Insert \$ FF" is its reduction by 1

Additional areas of memory may be allocated by the malloc indicated symbol. "Saw" only two types of Memory: Memory and the central bank of the extension, which resides in SPARTA.SYS procedure (bank system). The number of bytes that are be reserved more than an indicator of free memory, we in FAUX4 / 5 (\$ 0,785 / 6). To register X code representing the load type of memory (the so-called index memory: 0 Main, 2 bank system) and Y reset. With the additional allocation of memory, when a bank system lacks this place, the system tries to allocate main memory. When correctly performed the procedure (ie completed

positive), a pointer to the assigned area is in FAUX1 / 2 (\$ 0782 / 3), and index of memory - in the register X. Memory used by malloc then it can not be exempt, but this feature is intended for programs to reside permanently in RAM.

The allocation of memory banks

In addition to main memory and the bank system, an area which may TSR want to take, are free additional memory banks.

Allocation procedure is very similar to the above zademonstrowanego generates a list of sub-free banks.

```

; BANKALLOC
temp
= $ 80
LDY COMTAB + $ 1d
beq error
dey
Jan COMTAB + $ 1d
yeah
and # $ 03
asl
asl
Solid temp
yeah
LSR
LSR
tax
lda T_ + $ 08, x
now temp
eor PORTB
and COMTAB + $ 1C
eor PORTB
iny

```

clc
rts
error sec
rts

Subprogramme reserves (on a permanent basis, until the next cold start-up) a memory bank notes. When executed correctly (with C = 0) the battery pack will contain a code PORTB podłączający reserved bank in the computer XL / XE, while the register Y - appropriate value Axlon to register in the computer Atari 800

Access to the bank system

Access to additional memory, as it is written above, you will differently depending on its type. First, we'll simplify the issue access to the bank system.

NOTE: the area of RAM is conventionally called the bank the system can be located at different addresses, depending on the memory configuration selected by the user. In particular, need believe that the code Switching was outside the \$ 4000 -- \$ 7FFF!

Connecting the bank system is realized by calling EXT_ON the system (\$ 07F1) in argument transmitted in the accumulator. Disconnection together with the reintroduction the previous system of banks - not necessarily because originally it must be connected to main memory bank - provides a procedure EXT_OFF (\$ 07F4). For the latter not to pass on any arguments.

The argument for EXT_ON code is the additional memory, we wish connect. It can not be permanent, since after the first call at the time it is nowhere said that the additional memory at all, and after Second, even if, whether any portion of it was us assigned by DOS. *Therefore, the value is passed on to EXT_ON always have, in one way or another, to obtain the first from operating system.*

The number one method applies only to programs stored in

relokowalnym SpartaDOS format, in which at least one block Binary is to be automatically loaded into the extra memory. This memory is always a bank system, unless there is no place in it -- then block the program is loaded into main memory.

Code memory, which loaded the binary blocks intended to additional memory can be found at the time of its launch in variable indicated symbol EXTENDED. The program installs the rezydentnie in the system should remember it was then the value (after return to DOS in it is set at zero), then, during its action, pass it as an argument for the EXT_ON when you need to refer to the part residing in the bank system.

Method number two concerns TSRs, which need access to memory is not by himself but by other

System drivers. A classic example of such a procedure is **Ramdisk.sys**. The data may be requested from the program user, you should generally be saved for home memory. However, as well the program may be causing SPARTA.SYS procedure for requesting to read buffers DOS, and these may be in memory of the notes. In any case, Ramdisk.sys must therefore decide to which memory is perform record, or he read, either sector, or block status, or PERCOM block. Such a decision is taken on the basis of the indicated variable symbol SYSCALL. Each contains a memory index, to which the desired is reading, or with which the desired word, and then, if necessary, need, it is the value SYSCALL then be transmitted as EXT_ON argument to the call to connect the memory audience. Connection EXT_ON memory must have its counterpart in a later call EXT_OFF, where access to the bank

Page 22

the system will cease to be needed. Access to other banks enlargement As mentioned above, access to other memory banks realized by saving directly to the respective values registers controls memory, ie PORTB (\$ D301) in computers XL / XE and Axlon (\$ CFFF) in the 400/800 computers. As the value of the the record should take the results of procedures, or BANKALLOC FREELIST printed on the previous pages. Due to the Separate addressing ANTICA memory and CPU, when expressing PORTB good practice is to change only

the current value of those bits,
which are responsible for switching the
bank, and leave the rest without
changes. This is done in the following
manner:

```
; BANKSWITCH
lda PORTB
Solid old_pb
lda new_pb
eor PORTB
and COMTAB + $ 1c
; PBMASK
eor PORTB
Solid PORTB
```

To date, the value of PORTB is stored in
bajcie

labeled "old_pb", "new_pb" is taken from
FREELIST code

bank, which want to connect.

Switching program memory on one of the
banks of enlargement, when
will end the use of additional memory,
it should restore its
initial configuration. In the case of XL
/ XE is quite simple,
simply, as shown above, before remember
somewhere
PORTB value, and then change it back.

Page 23

With the extension Axlon matter is
complicated because the control register
banks only to write - so you can easily
switch banks in
one-way, but to restore the system
before the switch already not so easy
You can, because reading from \$ CFFF not
pay the state registry, but some
random garbage. Do not be so remember
your current configuration
memory before.

This difficulty can be resolved in two
ways. The first - and
easier - is blocking a program runs on
computers
other than XL / XE. This is achieved by
checking _800FLG
(COMTAB + \$ 1B). When switching memory

is critical for the speed of the action - as it is, for example, in the case of drivers CON64.SYS and CON80.SYS where to send the screen of each single character entails the memory banks - then after Just no other way out, how to reconcile with the fact that our code to 400/800 with a Axlon will not work. The second method involves the use of a trick. But the current banking system extension Axlon know DOS, it is (in principle only) on it switches. Remember we achieve this state by calling EXT_ON to the value of ensuring the provision of bank system (ie, EXTENDED or code returned by malloc in the registry X, see above). DOS switches the memory banks, but it is not us interested. When EXT_ON give control, you can enter your desired value to the registry Axlon control, which the bank will make available a program for in question. In contrast, the previous state is obtained simply by calling EXT_OFF.

Page 24

Chapter 4: Handling errors

The standard procedure for handling error (U_FAIL)

SpartaDOS handling error in the creation procedure

U_FAIL library with an error code in the accumulator. U_FAIL may be called explicitly to the user, but generally there

This is the secret, if this procedure, he finds some irregularities, it automatically produces the library system.

U_FAIL has the unpleasant feature that will never come back. In other words, calling the program is

interrupted and removed from memory,
open files are closed, and the system
shown on the appropriate console
the error message gives control to the
command interpreter. This is
Comfortable in the case of critical
errors, which occurred after
program, and so can not carry on,
however, relatively
often there is also a need to handle the
error within the program.
The creation of traps (U_SFALL)
Program and want to take over the
handling of a particular error may
zastawić
trap. Use the library U_SFALL procedure.
Before her
calling for the AX must registers load
properly and younger
Senior byte address space in the
program, which will be transferred to
control at the time the error occurred.
You can sign up for more than one trap -
reacts always the one that
was established recently. However, while
the number of traps set up
should not be prejudged, the program can
not have more than 10 at a time.
When an error occurs in the trap set up,
it will be above
all removed (this is a one-off). Then
loads the library
stack pointer value memorized by
U_SFALL, switches banks

Page 25

memory stored in a state there and
executes the stroke at the address JMP
indicated by the user in setting up the
last trap. Code
an error that occurred, is loaded into
the battery, a marker
N register markup processor - set to
one. The rest remains
unchanged, ie in particular do not close
the files that the program previously
opened.
Removing the trap (U_XFALL)

As a rzekło, the trap is removed automatically when there's mistake. However, if the error does not occur, the trap usually becomes unnecessary and should be removed "by hand". Use the procedure U_XFAIL. No pass it any parameters, it simply removes the trap set up recently. U_XFAIL does not alter the content of CPU registers. All the trap set up by the program are removed at the time its completion and commissioning of control in DOS.

Error message (U_ERROR)

U_ERROR displays on the screen system error message, whose code was passed in the accumulator. Before that is sent editor, character EOL (\$ 9B). This is part of standard maintenance procedures Boot error by U_FAIL.

An example of the use of traps
; setting pulapki to the address
over the vector trapptr, it contains the
address
; marked a label trap
lda trapptr
ldx trapptr 1

Page 26

JSR U_SFALL

; a procedure call system
, which is expected in the speeches of
error

...

; when the error was not, we remove
pulapke

JSR U_XFAIL

; when the error occurs, the system will
automatically put off
; pulapke and give control to that place
; with N = 1 and the code of error in
the accumulator

TRAP

bmi error

; here, the proceedings in the case

; the successful conduct of the
procedure

...

rts

; a service error

JMP error U_ERROR

Page 27

Chapter 5: Command-line treatment

LBUF and buffer index BUFOFF

Interpreterowi command issued commands

at the DOS, either directly

by a user with the keyboard, or read

from a batch file,

is saved in the buffer LBUF (line

buffer). It has 64 bytes and is

at COMTAB + \$ 3F. Current position in

the buffer, ie the

Generally the next parameter indicates

the index BUFOFF (buffer offset,

COMTAB + \$ 0A).

LBUF buffer BUFOFF index and provide

input for

library procedures dealing with command-

line treatment, and programs

normally do not have the necessary

attention to their content. May

however, sometimes have to read more

than once a

element of the command line. You'll need

to remember the state and BUFOFF

restore it before you call the library

(with

All of which is more BUFOFF setting it

on

the next item for processing).

It must be noted that the first

parameter is the LBUF

Generally caused the program name. At

the time of its launch

BUFOFF points out, however, for the next

parameter - as the name of the program

has already been read by the interpreter

in DOS commands. To download this

name, you must reset the BUFOFF.

Buffer COMFNAM

Exit procedures for processing the

command line, those at the very least,

which results in the form of text can be found at
COMTAB + \$ 21 This is a buffer COMFNAM
(complete file name) on
30 bytes in length. The associated
variable is the TRAILS (trailing

Page 28

space, COMTAB + \$ 1A), it contains
length in the
COMFNAM parameter.

While LBUF contains a command line, it
is COMFNAM
intended for a single parameter. LBUF is
just behind

COMFNAM, in principle, these two buffers
connect with each other. It follows from
this,

that a very long parameters (more than
30 characters) collected by the library
from the command line and wkopiowane to
COMFNAM can override
the initial part of LBUF.

Reading the command line

The command line consists of a group of
elements, ie, separated from
each file name spaces, switches,
options, etc. stored on
generally in the order, which is known
in advance and defined as syntax
the command. It follows from this that
the program to realize
reading the command line from beginning
to end as part of element
in most cases can anticipate what kind
of

parameter will be reached next. This
program does not need to parse
the command line to walk - though, if
the need arises,
of course, can - as the library offers a
procedure for processing
the most common types of parameters. We
will be discussing now so simple
cases where the parameter type is known
in advance, the question of analysis
more
complex command line leaving at the end.

The number (U_GETNUM)
Relatively simple to download and
interpretation of parameter
number. This may be a number between 0
and 65,535 given in
decimal or hexadecimal notation. The
call returns U_GETNUM
zero (Z = 1), where the parameter is not
a number. In wpadku (Z = 0)

Page 29

AX registers is appropriately junior and
senior byte values
binary equivalent to the number.

When the parameters are many more may be
separated by commas or
spaces. With U_GETNUM use a command such
as PEEK and Poke
SpartaDOS command interpreter.

The dwustanowy: ON and OFF (U_GONOFF)
Some commands you just something turns
on or off

giving respectively ON or OFF as a
parameter. For the treatment of this
type

parameter used procedure U_GONOFF
library. When the parameter
is ON, marker C register markup
processor is set to 1, and
If OFF - is at zero. When the parameter
is not specified neither he nor
OFF, the procedure reported error No 156
(Bad parameter) giving control
U_FAIL to the procedure and thus
disrupting the program. How to deal with
this problem, described in the chapter
titled "Using the errors".

In this way the command line overlay
KEY.COM proceed.

Options (U_SLASH)

The parameters for some of the programs
is conveniently provided in
a single "options" preceded by the
character slash "/".

To read such from the command line
procedure is used U_SLASH.

All options must be recognized place in
the table. One entry in this

array of options for one, two bytes, in
turn marker
of options and the corresponding letter.
For example, if
program wants to respond to the options
/ A / X, table should look like
as follows:
switch

Page 30

```
? a
. byte 0
. byte 'A'
? x
. byte 0
. byte 'X'
```

Address the array must pass the U_SLASH
in the records
AX, and the total size in bytes - in the
register Y. When the
analyzed the command line item is given
one of the expected
options, then the flag, in the above
example, they are
marked as "switch? a" and "switch? x ",
will take the value of \$ FF.
When given the option lacks the plate,
U_SLASH forward
U_FAIL control to the program by cutting
the error No 156 (Bad
parameter). In contrast, kept machined
when the parameter is not at all
option (ie does not start with a "/"),
U_SLASH returns to the place
Call not doing anything.
In this way the command line interprets
the MEM command.
Keyword (U_TOKEN)
Analysis of the command line for the
occurrence of certain words
provide key two procedures: U_GETPAR and
U_TOKEN.
U_GETPAR copies the next (ie those
indicated by BUFOFF)
parameter of LBUF to COMFNAM, BUFOFF
update, add length
parameter to the trails, and the vector
FILE_P - COMFNAM address.

U_TOKEN gets parameter text that is in COMFNAM

and tries to find it in the array of keywords, which address the transmitted in the registers AX.

Keywords must be placed in plate for a second, with the last character of each must be negatywie (ie setting bit 7). The end plate determinded by zero.

When U_TOKEN not match a keyword in the table, returns from

Page 31

skasowanym tag C. Otherwise, the finds marker C is set, a battery pack contains a sequence (ie token) of the keyword in the table, counting from zero.

In this way instruction is progressing, and a batch file IF SpartaDOS command interpreter.

Device Name (U_GETPAR / U_GEFINA)

When the parameter is the name of the device itself, for example, ID disk, you must download it to use steam procedures

U_GETPAR and U_GEFINA. U_GETPAR, as has been written above, gets the parameter of LBUF and insert it into COMFNAM, however, address

COMFNAM writes to the vector FILE_P.

U_GEFINA but gets parameter from the place indicated by

FILE_P, interprets it as the name of the device and, according to the

accordingly sets a record DEVICE (\$ 0761). His younger four bits

means the number of devices (such as the number of disk drives), while older

encode the type of equipment as follows:

\$ 0x - drive

\$ 1x - a clock

\$ 2x - a device CAR:

\$ 3x - a device CON:

\$ 4x - a device PRN:

\$ 5x - a device COM:

\$ 6x - a device NUL:

\$ 7x - reserved

When no ID is not specified, the code is adopted

keep the set device (generally, the current drive) downloaded from

Page 32

variable symbol CURDEV indicated. On the contrary, when given

ID is invalid and can not be decode, control

U_FAIL is transmitted to the error code number 130 (nonexistent

Device).

In the given way of U_GETPAR and

U_GEFINA uses such as command

CHKDSK.

The specification of the file (U_GETPAR / U_GEFINA)

Described above sequence U_GETPAR /

U_GEFINA may also be

used to get the file name. ID devices, as above,

is translated, then the value of DEVICE (\$ 0,761), while the rest of the specifications

is allocated to the path copied to the PATH (\$ 07A0, 64

bytes), and the name of the file-filled NAME (\$ 0,762, 11 bytes).

The name is translated into the format of an internal DOS,

ie to form NNNNNNNNXXX. When one part of the name, ie main or

extension, are less characters than the 8 or 3, is

supplemented by spaces, if any jokery (wildcards) are developed

in strings of question marks, etc.

This file is a specification for the

kernel SpartaDOS Strawn,

primarily for the driver DSK: in

SPARTA.SYS.

Calling it features a library system in general, the same exercise

U_GEFINA to jump in, so the user needs only to use

U_GETPAR before.
Specification directory (U_GETPAR /
U_GEPATH)
When the expected parameter is the
specification of the directory,
moving in the same way as above, with
the only change that, as with the other
procedures, instead of U_GEFINA, you
must call U_GEPATH. Review

Page 33

It translated the name of the device on
a code device (\$ 0,761) in the same way
U_GEFINA as a complete path is copied to
the PATH

(\$ 07A0, 64 bytes). NAME to nothing is
inserted.

This form of the specification file is
useful for calling
the kernel number 16 (kchdir) and 17
(kgetcwd). The relevant functions
libraries (CHDIR and getcwd) U_GEPATH
same cause, the program
user needs only so U_GETPAR call before.
The specification and file attributes
(U_GETATR)

Some commands, such as COPY or DUMP,
accepts as a parameter
specification file detailing the
optional attributes. For example,
DUMP + H FOO.BAR displays the contents
of a hidden file (with the attribute
+ H) FOO.BAR. Normally, however, a
hidden file will be ignored.

The task of analysis of such a parameter
U_GETATR performs the procedure.
He does it alone. ie there is no need to
call first

U_GETPAR. The battery must be
transmitted before the default
attributes for the file, if the user has
not submitted any. Mask
bit attributes compiled by default
following

Scheme:

+ \$ 01: protected (+ P)
+ \$ 02: hidden (+ H)
+ \$ 04: archive (A)

+ \$ 08: catalog (+ S)
+ \$ 10: not protected (P)
+ \$ 20: not hidden (H)
+ \$ 40: Do not archive (A)
+ \$ 80: no directory (S)

Page 34

Default masks are used most often: \$ 20 (not hidden) and \$ A0 (not hidden and not directory). U_GETATR returns to zero (Z = 1), whereas in LBUF is not enough parameters to be taken - such as given attribute, but not given names file. Mask attributes is inserted into FATR1 (\$ 0779), this is one of the criteria to seek the files in the office Library FFIRST and FNEXT (and thus also by all Other features of them benefiting from it, above all, fopen). Besides U_GETATR detail is the same as U_GETPAR, that is, as has been written above, it retrieves the parameter of LBUF and insert it into the COMFNAM, saves in parameter length trails, while the address COMFNAM entered in the vector FILE_P. Specification file with the default Mask (U_FSPEC) Some commands, such as COPY, accepts as a parameter path access with the specification of a file or a mask of selecting a group of files. This mask can be omitted, if it is to be '*' - COPY FOO> is to copy all of the directory (*.*) files FOO. Program executing the command must, of course, in the treatment posed parameters to identify whether a file name or mask was given, and when finds that it does not, add it. It is indeed simple, but onerous task - the city check several conditions - performs a function library U_FSPEC.

Called immediately after U_GETPAR or U_GETATR check whether located in COMFNAM parameter satisfies the conditions necessary for that add at the end of the mask *.* and, if so, add it whilst improving the value of trails. Combinations of different types of parameters

Page 35

The most common situation, when the program gets the first the name of the file, and then the options that begin with slash. Analysis such a command that is the difficulty, the program first should proceed as downloading a file specification (specific cases that are described above), and then call U_SLASH to read the options. Difficult case to the command of a kind in Pakistan, which has two form: ECHO ON / OFF or ECHO TEXT. In the first form enabled or disabled is performed by the echo command command interpreter. The second form simply displays the text on a given screen. The difficulty here lies in the fact that to identify the ON / OFF U_GONOFF need to call the procedure, and that if the parameter is neither an ON or OFF (but TEXT), permanently interrupt the program jumping to U_FAIL. The only solution is mortgaging the trap on the error, which may U_GONOFF generate. This has been described in the chapter titled "Handling errors".

Page 36

Chapter 6: the working name of the file Converting from 8 +3 format for internal (PRO_NAME) PRO_NAME library function converts the

file name
bufadr indicated by the vector (\$ 15)
with a Y-1 format
Internal and stores it in the buffer
NAME (\$ 0,762 - \$ 076C). In
Where the name is the name of the file,
the function returns the result of
non-zero (Z = 0) and a sign of EOL (\$
9B) in the accumulator. When it is
directory name, the result is zero (Z =
1), and the accumulator is
separator '>' or '<'.
Transforming the internal format
consists of splitting
File names for parts of central and
extension possible addition
both parts of spaces as possible and
develop Jokers
'*' In strings of question marks.
PRO_NAME is a service from which the
benefit described in
the previous chapter features U_GEFINA
and U_GEPATH. As a result,
including the user never normally have
no need to
directly cause.
Converting from an internal format for 8
+3 (U_EXPAND)
Reverse transformation, ie the name
stored in a format
internal and appearing in NAME (\$ 0,762
- \$ 076C) to a 8 +3
U_EXPAND function of the library. The
result will be entered into
buffer address AX, from the position Y.
The resulting string of characters ended
is a sign of EOL (\$ 9B).

Page 37

Chapter 7: environmental variables

What is the environment variable

SpartaDOS X is the only DOS on

ośmiobitowe Atari, which

implements of the larger computers known
environmental variables.

The environment variable is a string of
ASCII characters assigned

a unique name. These variables hold some

global settings
the system interpreter commands or
application programs.
An example of the first kind are the \$
PATH and \$ daytime,
\$ COPY second, third \$ MANPATH.
Environment variables are stored in the
additional memory
cache size of 256 bytes. The library
offers three functions
allowing for easy access to the data
contained therein.
Read variable by its name (getenv)
To read the variable name with a known
feature
Getenv. Name address the variable
function is found, must give
AX in (name should be completed sign EOL
mark
equality "="). Completion of a negative
result (N = 1) indicates that
buffer, this is not a variable. In the
opposite case, the result is
positive (N = 0), and the value of the
variable is enshrined in the form of a
string
ASCII character EOL culminating in the
buffer output package
mathematical LBUFF (\$ 0580).
Read by a variable number (NUMENV)
Alternatively, you can search for
variables not by name, but
according to their numbers further in
the buffer. We will do this procedure

Page 38

NUMENV. The number of variable need
provide in its battery pack when the
result
is positive (N = 0), the value of a
variable is recorded in the same
format and the same place as in the case
of getenv (see above).
This feature is generally used to read
all
environmental variables at a time, for
example, view them on
screen (as it does to the SET command

interpreter), search
whole, etc.
Record and erase the variables (putenv)
For the record variable in the buffer is
the putenv. New address
variable content to be given in AX. In
this address should be
find a string of ASCII characters end
with a EOL (\$ 9B) in the
the form: NAME = TEXT
This creates a variable NAME and assign
the text
"TEXT" as a value. When the buffer
environmental variable that
name already exists, is erasable before.
Comment on putenv same name variable is
a string
Ascii completed by EOL and free trade
equality
will remove the buffer variable with
that name.
NOTE: when combined with NUMENV and in
one putenv
loop in return for specific variables
and delete them, should be
remember that clears the variable
results in a reduction in numbers by 1
all the variables that are followed in
the buffer. Therefore, in
such a loop after calling putenv
kasującym variable DO NOT
increase the counter for NUMENV.

Page 39

Chapter 8: read and write files

Opening the files (fopen)

The device performs the function fopen.

Required parameters

Opening transmitted are as follows:

- 1) specification file should indicate
the vector FILE_P
- 2) the opening of the type opmode (\$
0,778)
- 3) mask attributes wanted to FATR1 (\$
0779)

The specification of the file indicated
by FILE_P should take the form of
text (ASCII string end with a EOL).

Specification

equipment must enter the convention SpartaDOS X, and not the Atari OS - ie eg A:> Catalog> textfile instead of D1:> Catalog> textfile.

Opmode has the same function as a byte ICAX1 Channel I / O XL OS when you open the file (when the device functions OPEN D: for opmode value through the CIO is charged with the ICAX1).

The value of this variable consists of two pólбайтów. Younger signals mode of access to the data:

\$ x4 - reading

\$ x8 - record

\$ x9 - annotations

\$ xC - data exchange (read / write)

Senior to mask, which has set subsequent bits

the following meanings:

7 - long format directory

Page 40

6 - attributes tracking mode

5 - opening up the route from the search (\$ path)

4 - direct access to the directory

Mask attributes wanted FATR1 (\$ 0779) is used at opening the file to read. The file name will be retrieved in the directory and open only when they meet the condition encoded

FATR1 bits:

+ \$ 01: protected (+ P)

+ \$ 02: hidden (+ H)

+ \$ 04: archive (A)

+ \$ 08: catalog (+ S)

+ \$ 10: not protected (P)

+ \$ 20: not hidden (H)

+ \$ 40: Do not archive (A)

+ \$ 80: no directory (S)

Normally when you open a regular file, the mask \$ A0 (Not hidden and not directory).

In the event of error control is transferred to U_FAIL (see Chapter on "Handling errors"). When the opening was successful, the file handle

(Eng handle) is entered into FHANDLE (\$ 0760).

Closing files (fclose / FCLOSEAL)

The library offers two functions here: fclose closes a specific file,

namely that the handle at the time of the call is in FHANDLE (\$ 0760). FCLOSEAL while closing all the files, which are open at the time.

Page 41

There is a possibility of security before closing the file FCLOSEAL. It is for this purpose to enter its grip FHANDLE (\$ 0760) and call the FCLEVEL with a value of \$ FF in the accumulator.

The same procedures for the value set on the battery the current value of SYSLEVEL abolishes protection.

Read and write individual bytes (FGETC / FPUTC)

Read and write a single byte file, whose handle is in FHANDLE (\$ 0760), exercise the functions FGETC and FPUTC. FGETC calls byte read in the accumulator. When the end there file, in the accumulator will mark EOL (\$ 9B), X value in the registry

\$ FF, a register of tags indicates the result will be negative (N = 1).

Any other error causes the transmission to control U_FAIL.

FGETC

does not change the contents of the register Y.

FPUTC sends a byte to a file over in the battery. Call

does not change the contents of the registers A, X and Y.

Speeches error causes

the transfer of control to U_FAIL.

I / O performed by FGETC and FPUTC for mobile devices

DSK: and CAR: mikrobuforowane by the library, so that

proceed much faster than the readings and records of

individual bytes

fread functions and fwrite.

Read and write records (fgets / FPUTS)

Read and write records file, which is in the grip FHANDLE

(\$ 0760) and carry out the functions fgets FPUTS. The

parameters for both

to be in the registers in the AX address a buffer overflow in Y its length.

The record is a string of ASCII characters end with a EOL and not

longer than 255 bytes.

Page 42

When read by fgets will correctly, the function returns from positive (N = 0), zero in the accumulator and the number of read

bytes in Y. When there is excess data, the result is negative (N = 1), and

battery pack is \$ FF. This situation represents a crash 137

(Truncated record) in the XL OS. Any other error results in automatic

the transfer of control to U_FAIL.

When the entry by FPUTS, if the string is shorter than the value of Y
Calling the time, must be completed sign EOL (\$ 9B).
An error causes an automatic call U_FAIL.
Fgets and FPUTS use of the functions of the library and FGETC
FPUTC, so for the device CAR: readings, and for DSK: records
and
readings results are mikrobuforowane.
Record formatted text file (FPRINTF)
Record formatted text file to implement the library
FPRINTF. It works identically to the one described in the next
chapter
function printf (these are two different entry to the same
functions), but
only that it is performed to save the file, which is in the
grip
FHANDLE (\$ 0760), and not on the console.
As in the case of record results, FPRINTF sends data
FPUTC through, so write to the disk is
mikrobuforowany.
Read and write blocks of binary (fread / fwrite)
Read and write binary blocks and carry out the functions fread
Fwrite. Parameters, in addition to the file handle FHANDLE (\$
0760)
pass with:
- FAUX1 / 2 (\$ 0782 - \$ 0783): address buffer

Page 43

- FAUX4 / 5 (\$ 0785 - \$ 0786): the amount of buffer
The size of the buffer must be greater than zero. In the case
of
the success of the functions of the result comes back positive
(N = 0). When the
Fread end of the file, the result is negative (N = 1). Any
other error
button to U_FAIL.
Fread and fwrite are not mikrobuforowane, because the use of
these
functions to read or write very small portion of the data
(after a few or
a dozen bytes) does not pay off. It is better for this purpose
and use FGETC
FPUTC.
Read the length of the file (FILELENG)
FILELENG read length feature an open file, which
handle is in FHANDLE (\$ 0760), and puts the result in FAUX1-3
(\$ 0,782 - \$ 0,784).
Changing positions in the file (FTELL / FSEEK)
To read the current position of reading or writing to the file
serves
FTELL function. It saves the records FAUX1-3 (\$ 0,782 - \$

0,784);

from the beginning of the file, the number of byte to be read or

saved in the next I / O.

The operation is the opposite, namely a change in this position carries out the function

FSEEK. New position have to give in FAUX1-3 (\$ 0,782 - \$ 0,784).

In both cases, handle the file must be saved in FHANDLE (\$ 0760).

When you try to set the position beyond the end of the file to open

Read control will be transferred to U_FAIL error No 166 (Range error). In contrast, a similar operation on a file open for writing

does not produce error, following that record data and closure

Page 44

will generally rise to the discontinuous (with the "hole", which are not assigned to any sectors of data).

Page 45

Chapter 9: console, entrance and exit

Record of individual characters on the screen (PUTC)

Record of individual characters on the screen (ie to the CON:

--

Here again, that the output CON: it can be user -

routed to any other file) performs a function library

PUTC. Mark it on to record a pass in the accumulator.

The call does not change the contents of the registers A, X, Y or variables

FHANDLE (\$ 0760) and DEVICE (\$ 0761).

Record record on the screen (PUTS)

Record record text to the CON: realizes function

Library PUTS. Address string pass in the library

AX registers, and Y serves its length. If the string is shorter than

Y value at the time of the call must be completed sign EOL (\$ 9B).

PUTS does not change the contents of the registers A, X, Y or variables

FHANDLE (\$ 0760) and DEVICE (\$ 0761).

Record formatted text on the screen (printf)

Printf function sends it to the console and the text strings of data

processed in the form of text and formatted according to the pattern. One of the features is that all data is transmitted by placing them directly for the upsurge in JSR printf.

Schematic call is as follows:

JSR printf

Page 46

. byte pattern formatujący, 0

any data

...

"The pattern formatujący" is a string of a length to 253 characters ending zero. In the simplest case, the plain text to view, for example:

JSR printf

. byte "Sugar krzepi", 0

A window is simply a given text on the screen, stopping the cursor

at the end. In order to connect with the transition to the new line, you need to

the end of the string (before zero) add character EOL (\$ 9B).

But the whole strength of the printf function lies in the fact that given the

characters, which is, again, only the model format text

output, you can put formatting commands. Here is a list of them:

%% - Write a single character%

% c - write a single character that is located at the following address

% s - write a string of address

% p - the same, only instead of the address is the address of its index

% x - value from the specified address, write as many 16-bit hex

% b - value from the specified address, write to be 8-bit number dec

% d - the same as the only 16-bit number dec.

% e - the same value of 24-bit

% l - the same, the 32-bit (from SpartaDOS X 4.40)

One command formatujacemu must correspond to separate

placed as a benchmark indicator format (in the above

the pattern of these indicators have been named any data).

Page 47

Using this is very simple. Suppose that the label symbolizes value

address 16-bit words stored in the usual convention

junior / senior. We want to transform this value to strings

representing the number in hexadecimal and decimal:

JSR printf

. byte "VALUE = \$% x =% d", \$ 9B, 0

. word value, value

When we value at the value of \$ 98AB, appears on the screen

We think:

VALUE = \$ 39,083 = 98AB

Wrote above that% x is to format the "16 --

bit "hexadecimal numbers. That is indeed, but only if

zażyczy programmer does not itself differently (ie,% x defaultowo treats values as 16-bit). In fact, the library reads All numeric values as 32-bit (or, in versions SpartaDOS older than 4.40 - 24-bit). When the value is less than or greater defaultowe than 16 bits, we need to provide, for example, JSR printf

```
. byte "VALUE = $% 2x", $ 9B, 0
. word value
```

take into account only the youngest byte value at the address value and bring it as a hexadecimal number dwuznakową. By contrast:

```
JSR printf
. byte "VALUE = $% 8x", $ 9B, 0
. word value
```

Page 48

Doing so that the number printed on the screen will have all eight digits. The number specified here as a sign of% must be a number of decimal range from 1 to 255 (in fact, it may be up to 65535, but the older byte will be ignored).

In all the examples given, when a string of digits to print starts with zeroes, they will be truncated. However, you can enforce their

removal preceded by zero value, meaning the the expected length of the string, for example,

```
JSR printf
. byte "VALUE = $% 04x", $ 9B, 0
. word value
```

In the end, it (ie, the length of the string, in the most recent example, "4") does not must be constant, can induce a library to retrieve it with a variable in

this way:

```
JSR printf
. byte "VALUE = $% 0 * x", $ 9B, 0
. numlen word, value
```

The label indicates numlen place in memory where they should printf

get the desired length of the digits.

Similarly, formatting strings decimal digits: giving the target number of digits for wildcard% b% d% l, we can not truncate

derived string of characters to their desired number.

In the case% c derived is always one character, all Additional attributes formatting commands are ignored.

% s removal causes the console ASCII string

located at the following address. This string should be

ended zero. At the% may show the number of characters. Then,

when

Page 49

string is longer, it will be truncated, and when it is shorter, it will be padded with spaces. %p works identically to %s the only place that address the address given dwubajowego index containing the address. Printf does not change the contents of the registers A, X, Y or variables

FHANDLE (\$ 0760) and DEVICE (\$ 0761).

Read a single character to the console (GETC)

The task of reading a single character to the console meets function

GETC. Read byte is returned in the accumulator. When reading to be successful, the result is positive (N = 0), and X is loaded

\$ 00 When there end, it comes back negative

(N = 1) and the value of \$ FF in the register X. Any other error causes

automatic transmission control to U_FAIL (see Chapter "Handling errors").

When you need a byte is not reading from the console, but the keyboard

should be used U_GETKEY function. It is waiting for the press key, then returns the ASCII code in the accumulator. Remember However, the latter must, that it exists only in versions of the X SpartaDOS

4.40 and above.

Read the record from the console (GETS)

Read the record from the console carry out the function gets.

The record is

a string of ASCII characters end with a EOL (\$ 9B). As for PUTS parameters to be in the registers in the AX address a buffer overflow in Y

length. Records can not be longer than 255 bytes.

When the transition is read correctly, the function returns the result of

positive (N = 0), zero in the accumulator and the number of bytes read in

Y. When there is excess data, the result is negative (N = 1), and

Page 50

battery pack is \$ FF. This situation represents a crash 137 (Truncated record) in the XL OS. Any other error results in automatic

the transfer of control to U_FAIL (see the "Handling errors").

Features PUTC, PUTS, printf, and GETC GETS performing reading and

record data for the console are just a separate entry (the so-

called wrappers) to described in the previous chapter, general features and write read files FPUTC, FPUTS, FPRINTF, FGETC and fgets. Using the former rather than the latter allows painlessly passed on to

library system to handle redirects I / O.

Redirection of I / O (DIVIO / XDIVIO)

Entry and exit from the console can be redirected to any other file by using the DIVIO. She is the

SpartaDOS command interpreter in a situation where the user for example, command

DIR>> file

.

DIVIO requires a vector FILE_P specification for the file to which (or whose) is to be established with the console redirection

(or her). The register Y included \$ 00, where is redirected output (ie PUTC features, etc.), a \$ 01, where the entrance is to be rerouted

(GETC). Possible errors in the opening are reported to redirect

U_FAIL. In the case of the success of your grip, where it is redirected data will be in FHANDLE (\$ 0760).

The task of the opposite, this is the end of the redirection, pursues

XDIVIO function. Parameter serves as \$ 00 or \$ 01 in the register Y, so

Same as above. Handle the file should be redirected be FHANDLE (\$ 0760). The call will close XDIVIO this file.

When entering the console has been redirected and end of the file,

XDIVIO library calls automatically. The same thing with

Page 51

closing all files and after the onset FCLOSEAL error.

Vector output (PUT_V / VPRINTF)

Another method of transfer, this time the only exit for the console, is

PUT_V defense of the vector. Normally he is not any meaningful value, the program must set the address of the servicer record. It may be, for example, the procedure for making direct entries to

memory screen, with the exception of equipment CON: or E:. It should be

by the end RTS, it is desirable to store the value of records.

When the file handle in FHANDLE (\$ 0760) is set to 100 (\$ 64), FPUTC call button indirect (JMP) by PUT_V with the to the view set out in the accumulator.

Since FPUTC is implicitly called by FPUTS and

FPRINTF, so the output of these functions is thereby also relayed.
FPRINTF Instead, it is better to do so by using VPRINTF - is Another (third now) entrance to the same procedure, and its action is exactly the same as printf, with the only difference being that the active channel I / O is automatically switched to handle 100 before text, and also automatically switched back after him. A programmer does not have to be so concerned about the content of FHANDLE (\$ 0760) in this case.

Page 52

Chapter 10: Catalogs

Search files (FFIRST / FNEXT)

To search for files in directories indicated designed two Features: FFIRST and FNEXT. Parameter input is the first complete catalog with specifications given at the end of a mask files

FILE_P indicated by the vector, and attributes sought in FATR1 (\$ 0779). Found the entry is placed in DIRBUF (\$ 0,789 - \$ 79F) in the

the form of "raw" entry in the directory format SpartaDOS (ie such as is saved to disk, see the "X Handbook SpartaDOS user ", Chapter 7).

FNEXT not require any additional parameters, searches simply following the entry criteria set by the previous function

FFIRST.

Signaled the end of the catalog is in both function by the return of

a negative result (N = 1). Any other error button to U_FAIL.

NOTE: FFIRST in fact open to specify a directory

read a file handle is placed in FHANDLE (\$ 0760). Therefore, after

finished searching, it is necessary to call for fclose the handle in order to close the file directory.

Read the catalog (FDOPEN / FDGETC / FDCLOSE)

The library includes three features which provide a catalog of formatted, which is legible for humans (such as such, what appears

on the screen after the interpreterowi commands command DIR). They are

turn FDOPEN, FDGETC and FDCLOSE. FDOPEN open stream

data directory, FDGETC collected from him asking for it in a byte

accumulator (and showing a positive result in the registry Tag

Page 53

CPU), and FDCLOSE directory open previously closed by

FDOPEN. FDGETC errors reported to U_FAIL, with the exception of status the end of the file - when takowy occurs, the function returns to the program calling from a negative result (N = 1). Each of them does, in principle, only one: call the appropriate misc_entry feature, namely the order No. 6 (misc_fdopen), 7 (misc_fdgetc) and 8 (misc_fdclose). Read formatted catalog so you can pursue by the way (ie by the vector misc_), where the library system is unavailable due to the launching of the command X. You just need to remember that misc_ never gives control to U_FAIL, paying Just after the error code in the battery pack instead of a further byte read by misc_fdgetc not be provided in the registers, but on Page zero ICAX6Z (address \$ 2F). The opening is a specification for FDOPEN directory FILE_P determined by the desired manner and format in Opmode (\$ 0,778). SpartaDOS X versions older than 4.41 are familiar with only two how to format: The format of "long" SpartaDOS (opmode = \$ 80) and format "short" AtariDOS (opmode = \$ 00). In SpartaDOS from version 4.41 and above is considered by the opmode FDOPEN as a mask selecting bits (separately) the means formatting the different elements of the directory. The importance of bits:

- + \$ 80 - long format directory
- + \$ 40 - display the attributes
- + \$ 20 - insert a space between the name and extension
- + \$ 10 - full stop after the name (if bit 5 = 1)
- + \$ 08 - two spaces before the name, '*' for a secure file
- + \$ 04 - the time in the 24-hour format
- + \$ 02 - when the long format directory (Bit 7 = 1), without a second time in short format will be displayed name extensions directory

Page 54

(instead of the standard

f ... "

), The directory will be denoted by a colon before the name.

+ \$ 01 - display the size of the sectors in (the bit 7 = 0).

To maintain compatibility with older versions SpartaDOS X given by the user is transmitted to \$ 0B, and \$ 80 to \$ A2.

NOTE: misc_fdopen, and thus also function libraries

FDOPEN, the discreet opening the file directory to read. In Therefore, (and) should not be neglected FDCLOSE call (or misc_fdclose), when the reading is completed, and (b) may be open only one file at a time, because misc_ is unable to remember more handles.

Page 55

Chapter 11: loading of binary

Load the program into memory (U_LOAD)

U_LOAD symbol indicates SpartaDOS binary loader. It fills three tasks: (a) the loading of binary programs into memory along with mobilization, (b) without loading run, and (c) run them.

Inputs U_LOAD is a specification file indicated FILE_P vector control and value in the registry FLAG.

Indicated

the file is opened for read-compatibility mode with \$ PATH (Opmode = \$ 24) and, if this is correct binary is loaded into memory. Any errors cause interruption of the operations and jump to

U_FAIL.

FLAG register can decide what to do with the loaded program. When

7 FLAG bit is zero, the program is started immediately. In Otherwise, only the updates are free memory (for programs in the format relokowalnym SpartaDOS), and

control returns to the caller. Mobilization (already without rechargeable) follows the re-call of U_LOAD the same specification file and registry skasowanym bit FLAG 7.

To remove the program from memory (U_UNLOAD)

When the program was running normally, its removal from memory The following soon as the interpreter to give control commands SpartaDOS. In the opposite case (loaded with FLAG> \$ 7F) in the

to remove the code from memory must call the U_UNLOAD.

It removes everything that the user previously load by U_LOAD.

Page 56

Chapter 12: File Management Functions

Change the file name (RENAME)

Change the name of that file carry out the function RENAME. Parameter input specification is complete a file name which is to be changed. These specifications should indicate the vector

FILE_P. After the source specifications should be, separated a space or a comma, name, which is to be attributed to the

file. Possible

Errors are reported to U_FAIL.

In earlier versions SpartaDOS X can be, using RENAME, several files in the directory to give this one same name. From version 4.40 is impossible.

Deleting a file (REMOVE)

Erases the file performs the function REMOVE. Parameter input specification is to remove the file indicated by FILE_P. An error button to U_FAIL.

Removing directory (RMDIR)

Erases the directory function RMDIR implement. Parameter input is the specification of the directory (without final separator)

indicated by FILE_P. An error button to U_FAIL.

The catalog must be empty. Removing the directory together with those in

the files and folders to other opportunities outside the library. When

necessary, call contained in the module ROM program

DELTREE.COM with the name of the directory to delete it as a parameter.

Page 57

The creation of a directory (MKDIR)

Creates a new directory feature MKDIR. Parameter input is specification of the directory (without final separator) indicated by

FILE_P. An error button to U_FAIL.

Changing the attributes (chmod)

Change the attributes of the file performs the function of chmod. Required

parameters are determined by specification file FILE_P, mask wanted attributes in FATR1 (\$ 0779) and the new mask attributes in FATR2 (\$ 077A).

Mask wanted attributes elect files, which attributes to be amended. It works the same way as the other taking into account the functions (such as fopen):

- + \$ 01: protected (+ P)
- + \$ 02: hidden (+ H)
- + \$ 04: archive (A)
- + \$ 08: catalog (+ S)
- + \$ 10: not protected (P)
- + \$ 20: not hidden (H)
- + \$ 40: Do not archive (A)
- + \$ 80: no directory (S)

In the mask of attributes defined individual bits have the following

the importance (when set to 1):

- + \$ 01: security (+ P)
-

Page 58

+ \$ 10: unprotection (P)
+ \$ 02: hide (+ H)
+ \$ 20: disclosure (H)
+ \$ 04: archive (A)
+ \$ 40: niearchiwalny (A)

Attribute S (directory), naturally, there is no way to change this.

Error gives control to U_FAIL.

Make sure you download the BOOT (SETBOOT)

SETBOOT sets indicated (by FILE_P) as the binary file, at startup, which is to be automatically loaded by loader located at the beginning of the drive. This function of course

only on drives formatted SpartaDOS. The possible causes confusion

the transfer of control to U_FAIL.

Page 59

Chapter 13: Other features Disc

Formatting the drive (format)

FORMAT function does not do anything, other than to call on the screen

SpartaDOS Formatter menu. The caller did not provide it no parameters.

In earlier versions of X SpartaDOS call by the formatter require prior "manual" (by interference in the records banks transfer module) cartridge's a switch to the bank 0 From DOS version 4.40 at this inconvenience has been removed, the system itself

selects the appropriate bank, to the formatter no longer resides in a bank 0

However, if we want to build compatibility, it is better to use the XIO 254

Record fresh directory (builddir)

Builddir takes a "soft" format, it's just new blank saves from the root, bitmap disk and bootsektor SpartaDOS format. It is the only method of formatting ramdisków

and partition the hard disk. Input parameters are:

- 1) code device in DDEVIC (\$ 0300). \$ 31 for the drive
- 2) the number of devices in DUNIT (\$ 0301).
- 3) a pointer to the new label disc in AX.

The label is eight characters ASCII. If he were to be shorter, you must

supplement it with spaces to that length. The label can not start with

sign space.

The function automatically reads the configuration of the drive and on this

determining the rest of parameters (number of sectors, etc.).

In SpartaDOS

X 4.40 is automatically turned on optimizing bitmap (function Optimize formatter). The status of the operation is returned to the registry Y.

Page 60

Can not read the disk parameters (GETDFREE)

GETDFREE read different information on the disk and puts them in

PATH (\$ 07A0-\$ 07DF). Before you call have to give specifications

equipment indicated it FILE_P vector. Any errors are reported to U_FAIL.

Data (17 bytes) are placed in the PATH as follows:

- + \$ 00: the identification code file
- + \$ 01: the number of encoded bytes in the sector
- + \$ 02: the total number of sectors (2 bytes)
- + \$ 04: current number of free sectors (2 bytes)
- + \$ 06: the name of the disk (8 bytes)
- + \$ 0E: the number of sequential
- + \$ 0F: the number of random
- + \$ 10: byte irrelevant

Identification code file to see what system

Files can be found on the disk, as follows:

- \$ 0x - filesystem AtariDOS vx0 (eg \$ 02 = DOS 2.0)
- \$ 11 - disk file SpartaDOS v. 1.1
- \$ 2x - disk file SpartaDOS v. 2.x
- \$ 3C - PC-MIRROR
- \$ 40 - filesystem module SpartaDOS (CAR:)
- \$ FF - disk file MyDOS

NOTE: GETDFREE will call the kernel No \$ 13 From version 4.41 SpartaDOS format of the data returned by the kernel call differs completely from what draws GETDFREE. In SpartaDOS 4.2x it was not so - the kernel \$ 13 in the 4.4x is incompatible backward-compatible with 4.2x. For the sake of compatibility programs should be

Page 61

refer to it only by the library (ie by GETDFREE), and in the absence of such a possibility, the system function XIO 47

operational. Both of them to convert data returned by the kernel

to the "old" format, in accordance with 4.2x

Change the current directory (CHDIR)

CHDIR changes the current directory on the drive, which specification

FILE_P vector points. Error causes U_FAIL to jump.

Read the current directory (getcwd)

Getcwd reads the path from the root of the drive to

current directory and puts it in your PATH (\$ 07A0-\$ 07DF).
Specification of the disk (as text), indicate the vector
FILE_P. Error results in calling U_FAIL.
In earlier versions of the X SpartaDOS getcwd not pay
name extensions directory. From version 4.40 has been
corrected.

Page 62

Chapter 14: ancillary functions

32-bit multiplication and division (MUL_32/DIV_32)

The library contains two procedures for calculating the total
figures

Free trade: MUL_32 performs the multiplication of two numbers,
32-bit, and

DIV_32, respectively, sharing none.

Ingredients operations must be entered in the order of bytes
the reverse of the normal (oldest first byte!) as follows:

1) multiplicand (or dividend) under COMTAB + \$ FF (4 bytes)

2) the multiplier (or divisor) under COMTAB + \$ 0103 (4 bytes)

The result of the calculation, in the same way in reverse
order of bytes

We take from COMTAB + \$ 0107. Both the signal

overflow by setting bit register markers C (C = 1).

Replacing the small letters on the large (ToUpper)

The call ToUpper code ASCII lower-case letter in the
accumulator

turns it into a large letter code. If the code does not mean
small

letters, it is not done nothing whatsoever.

Checking the directory separator (CKSPEC)

The procedure CKSPEC verify that the mark passed in the
accumulator

is one of the following: ':', '>', '\\', '<' (separators are
likely to

occur in the specification file), and if so, returns with the
result

zero (Z = 1). In the opposite case, the result is non-(Z = 0).

Page 63

Chapter 15: procedures for initiating

Initiating overlays after RESET (S_ADDIZ)

Part overlays must be initiated after each

Pressing button. For example, it can be all kinds of drivers

Install the CIO, as contained in SpartaDOS X 4.41 CON64.SYS

and CON80.SYS. When you reset the system they need at least
re-install the OS array handlerów-in.

For the recording of such overlays is the S_ADDIZ library.

Please pass to it in the registers AX address procedures for
initiating

overlay. DOS entered this address to a special queue, and
after RESET

performs in turn sub-registered there.

A queue is only five seats, the lack of registration is signaled by the return from S_ADDIZ with an tag C Tag Registry (C = 1).

Exit to DOS-u (_DOS)

Calling JMP _DOS has the same effect as JMP (DOSVEC)

- The caller is finished and removed from memory, and control returns to the command interpreter. A better way to the end

program is to implement an order to RTS.

Warm restart SpartaDOS (_INITZ)

Calling _INITZ causes warm restart SpartaDOS. All the files are closed, the settings reset, resident drivers kernel initiated from the beginning.

_INITZ Is one of the procedures triggered after hitting

Page 64

Reset by the user. In earlier versions SpartaDOS X execution resulted _INITZ the current paths of all drives on the main directories. From version 4.40 driver SPARTA.SYS

distinguish the call origination from the measure _INITZ from each other and

paths are not reset.

Page 65

Chapter 16: the manipulation of symbols list

Searching the list of symbols (S_LOOKUP)

The system contains the procedure for searching the list of symbols, and

add and remove symbols. Normally, however, these programs do not

benefit from them directly, because all matters related to deals with the symbols of binary loader SpartaDOS. List of symbols

attached to blocks of binary program is, after loading the code to

memory, it automatically translated into addresses, and these are inserted in

adequate space program. Similarly the case definition

new symbol of the loaded program, this definition encoded

is the same binary structure, and creating new

symbol and supplying it to the list of global deals with the loader.

Sometimes, however, there is a need to "manual" search of the list

symbols by the program. For example, is located in the ROM disk CAR:

Z. SYS driver needs access to a symbol I_FMTTD

defined by TD.COM, but can not be dependent on its

presence, because then it would not be possible without the

load Z. SYS
prior to load TD.COM (inability "zresolwowania" symbol included on the list annexed to interrupt the loading and "Loader: Symbol not defined").
In this context, you must use a procedure S_LOOKUP.
It requires registration of the name of the symbol, as supplemented by up to eight characters spaces, if necessary, under the symbol + \$ 02 Calling JSR S_LOOKUP
Calls to zero (Z = 1), where the wanted symbol does not exist. Otherwise
Otherwise, the result is different from zero (Z = 0), and under the symbol + \$ 0B and SYMBOL + \$ 0C sequentially is a junior and senior byte address reported by a symbol, but we have EXTENDED index memory. EXTENDED byte is well before reset.

Page 66

Since the same procedure S_LOOKUP, structure and characters EXTENDED variable symbols are indicated, it is clear that this path access to the list of symbols is open only for programs written in SpartaDOS binary format. Besides simple binary Atari DOS-in usually do not need access to the list of symbols. However, if there had been such a need, there is a second entrance to S_LOOKUP, not quite, that are under constant address, it's a lot more simpler to use. This procedure, called the FSYMBOL, introduced in SpartaDOS v. 4.40. Entry to it is at \$ 07EB. An input address is supplemented by space-name symbol over in the registers AX (mł. / st.). In the event of a lack of symbol procedure returns to zero (Z = 1), otherwise the records will be AX contain the address pointed to a register Y - an index of memory.
NOTE: at the time of the launching of the X command list of symbols and the procedures described here will become available to you.

The addition symbol (S_ADD)
After calling S_ADD symbol indicated parameters will be included on the global list of symbols. Data entry pass in SYMBOL structure, it seems, for the record, as follows:
+ \$ 00 - \$ 01 rate for the next symbol (2 bytes)
+ \$ 02 - \$ 09: the name of the symbol (8 ASCII characters)
+ \$ 0A: byte control
+ \$ 0B-\$ 0C: address indicated by the symbol (2 bytes)
S_ADD requires the fulfillment of calling names (with

complementary spaces to eight characters, if shorter) and the address in bytes SYMBOL + \$ 0B and the + \$ 0C. Index

Page 67

memory should be recorded in the extended, as long as there is no longer there.

Byte control is completed automatically, just a pointer to symbol next on the list.

The symbol is added in the place indicated by MEMLO,

This rate is then increased by 13 bytes.

Stripping (S_CLEAR)

S_CLEAR procedure removes the global list of the symbols of symbols,

where PID is greater than keep the set. It is fundamentally a part of

U_UNLOAD procedures, and "keep your 'number of applications may not

be explicitly changed by the user. Therefore, the objective S_CLEAR export to the global list of symbols is not quite clear.

Page 68

Chapter 17: other symbols library

In addition to the above in the list of symbols, there is still a handful so far

not covered, and some were not even mentioned. These are the variables:

STATUS, SYSLEVEL, H_FENCE, tables CARVARS, DEVSPEC,

Devname, COMTAB2, procedures _CIO, _edit, and _CRUNCH

XDFREE. Most of them applies only to special

system drivers, a description of their functions and their use will be in

future in a supplement to this paper on write drivers.

XDFREE symbol indicates an extensive procedure to replace

GETDFREE, but since it was introduced to SpartaDOS

X 4.41 so fresh that, in principle, it can be described as experimental, its

description is omitted for the time being.

Page 69

Index procedures and system variables

_CIO

68

_CRUNCH

68

_DOS

63

_edit

68

_INITZ

63
BUFOFF
27
Builddir
59
CARVARS
68
CHDIR
61
Chmod
57
CKSPEC
62
COMFNAM
27
COMTAB
14
COMTAB2
68
CURDEV
32
Devname
68
DEVSPEC
68
DIV_32
62
DIVIO
50
EXT_OFF
20, 23
EXT_ON
20, 23
EXTENDED
21, 65, 67
FCLEVEL
41
Fclose
40
FCLOSEAL
40
FDCLOSE
52
FDGETC
52
FDOPEN
52
FFIRST
52
FGETC
41

Fgets
41
FILE_P
39
FILELENG
43
FLAG
55
FNEXT
52
Fopen
39
FORMAT
59
FPRINTF
42, 51
FPUTC
41, 51
FPUTS
41, 51
Fread
42
FSEEK
43
FSYMBOL
66
FTELL
43
Fwrite
42
GETC
49
Getcwd
61
GETDFREE
60
Getenv
37
GETS
49
H_FENCE
18, 68
I_FMTTD
65
INSTALL
18
LBUF
27
Malloc
18
MKDIR

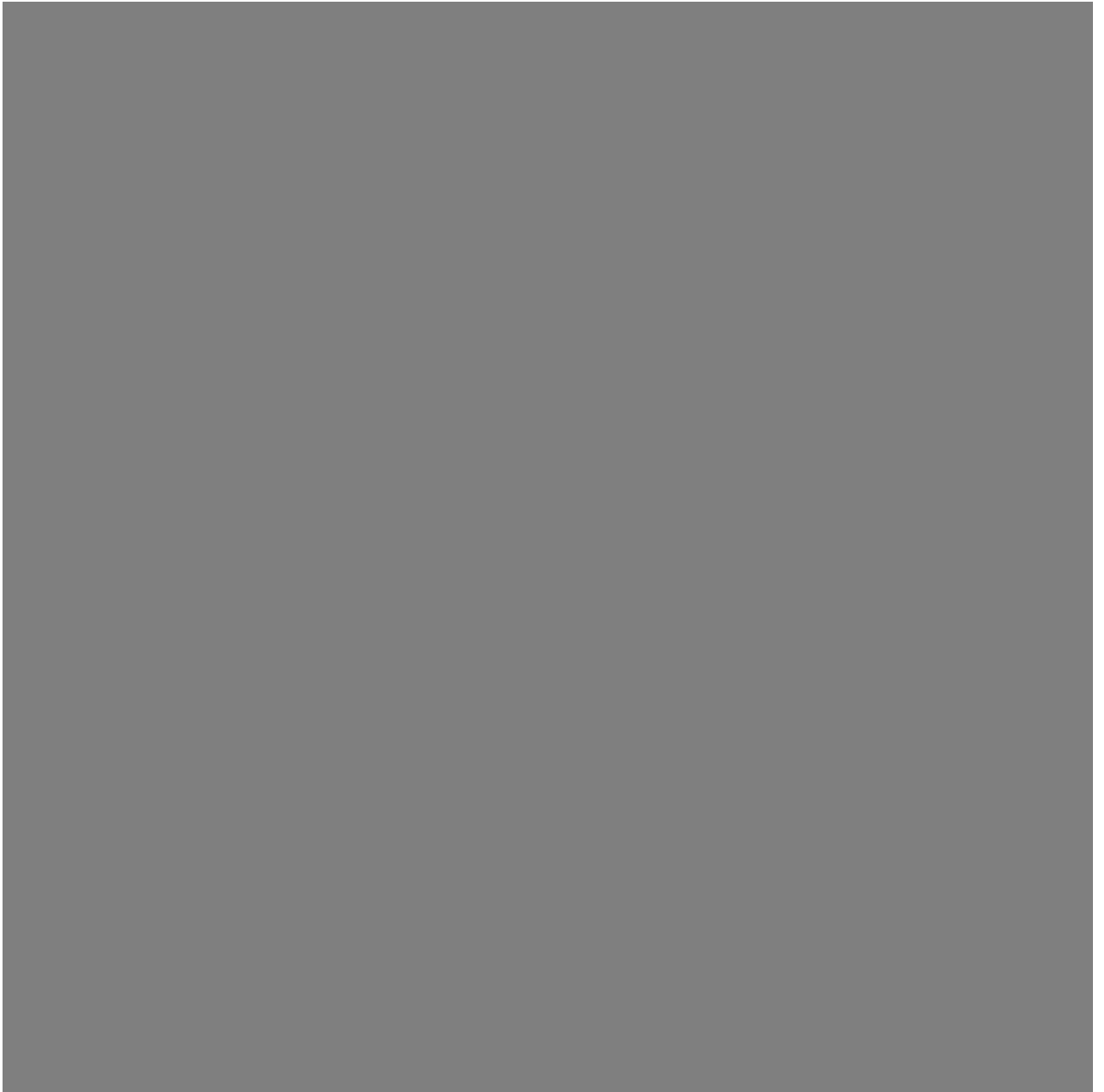
57
MUL_32
62
NUMENV
38
Printf
45
PRO_NAME
36
PUT_V
51
PUTC
45
Putenv
38
PUTS
45
REMOVE
56
RENAME
56
RMDIR
56
S_ADD
66
S_ADDIZ
63
S_CLEAR
67
S_LOOKUP
65
SETBOOT
58
STATUS
68
SYMBOL
6, 66
SYSCALL
21
SYSLEVEL
41, 68
T
15, 16
ToUpper
62
U_ERROR
25
U_EXPAND
36
U_FAIL
24

U_FSPEC
34
U_GEFINA
31, 32
U_GEPATH
32

Page 70
U_GETATR
33
U_GETKEY
49
U_GETNUM
28
U_GONOFF
29, 35
U_LOAD
55
U_SFAIL
24
U_SLASH
29, 35
U_TOKEN
30
U_UNLOAD
55
U_XFAIL
25
VPRINTF
51
XDFREE
68
XDIVIO
50







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Original Polish text:

je sobie do ulubionego asemblera.

[+ Suggest a better translation](#)